

Design Principles for Mobile Computing and Communication Systems

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Introduction and Overview

- Nokia Research overview
- Networking moving to distributed control
 - System design moving from centralized to distributed
 - Reduce cost of operation
 - Emerging pervasive networks
- 4 design paradigms for self-organization
 - Based on analysis of existing protocols
- Two examples in multi-hop wireless networks
 - Distributed resource control in multi-hop networks
 - Gateway discovery and multi-hop handover
- New challenges for emerging ubiquitous networks
 - From distributed to self-organized system design



Main reference: Christian Prehofer, Christian Bettstetter, **Self-Organization in Communication Networks: Principles and Design Paradigms**, IEEE Communications Magazine, July 2005. **NOKIA** Connecting People

Definitions of Self-organization

Merriam-Webster:

- 1: organization of oneself or itself
- 2: the act or process of forming or joining an organization. (as a labor union).

Wikipedia: Self-organization is a process in which the internal organization of a system, normally an open system, increases in complexity without being guided or managed by an outside source. Self-organizing systems typically (though not always) display emergent behavior.

Chemistry: The capability of a system to spontaneously generate a well-defined supramolecular assembly by employing its own components in a given set of conditions

Economics: A market economy is sometimes said to be self-organizing.

"We define a self-organizing system as one where a collection of units coordinate with each other to form a system that adapts to achieve a goal more efficiently."

Self-X

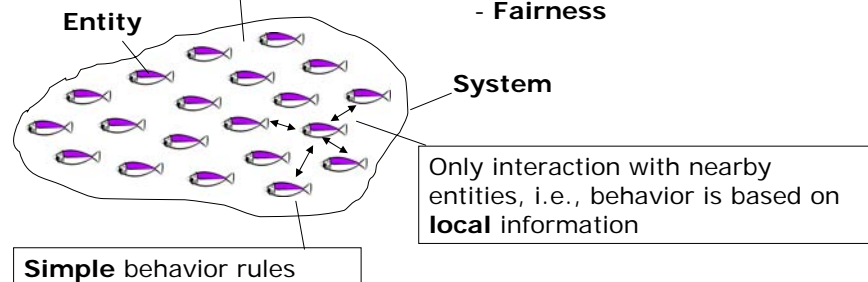
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Self-organization in Networking

Distributed control, i.e., no central or external control

Create system-wide properties

- Order
- Adaptability
- Robustness
- Fairness



- Need to understand typical „design patterns“ of self-organization
- Need to understand scope – which functions to be self-organized

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Summary of Paradigms and Examples

Paradigms

1. Achieve global properties with local behavior
2. Do not aim for perfect coordination
 - Implicit coordination
3. Minimize long-lived state
4. Adaptation

Examples

- Addressing and Naming
 - IPv6, ad-hoc networks
- Resource and congestion control and access
 - TCP
 - Medium Access (MAC) protocols
 - **Distributed resource control in multi-hop networks**
- Resource discovery
 - **Proactive and reactive discovery**
 - **Multi-hop handover**
- Data collection in pervasive networks

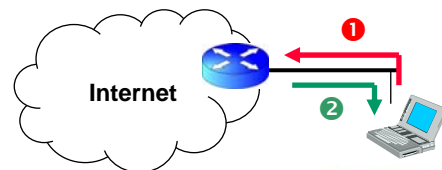
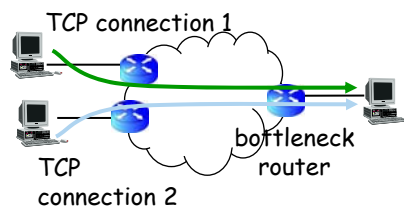
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Paradigm #1: Achieve Global Properties with Local Behavior

- Behavior rules based on local environment
 - Don't need to know about global state
- No need for central control
- System-wide property from local behavior

Examples

- **TCP protocol** for Internet congestion control
 - TCP is fully distributed, very simple, flexible
 - End-points control resources in the network
 - TCP achieves „fair“ resource distribution
- **Addressing** in the Internet – IPv6
 - Stateless autoconfiguration
 - Assign a globally unique address
 - Hosts configure address themselves
 - Global Address = Network prefix + Interface ID



Router Solicitation and Router Advertisement with Network prefix

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From Local Behavior to Self-organization

Some typical questions

- Local behavior and control are important – but does it work?
- When does local behavior lead to optimal solutions?
 - But optimal solutions are difficult in dynamic networks
- Local behavior rules may lead to inconsistencies
 - When is this preferable or acceptable?



Example for addressing

- **Passive autoconfiguration** for ad-hoc networks
 - Assign addresses locally based on local information
 - Maintain list of assigned addresses locally
 - Passively collect information about already assigned & used addresses
 - Detect address conflicts in a timely manner locally



Paradigm #2: Do not Aim for Perfect Coordination: Implicit Coordination

„Talking is silver, silence is golden“

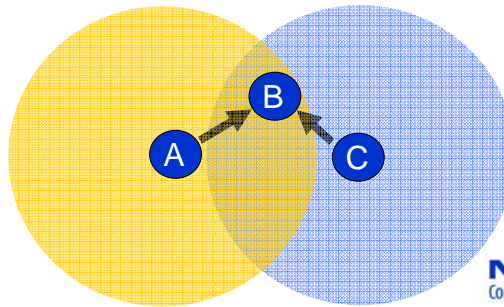
- Robustness often more important than perfect or optimal solution
- When can we admit inconsistencies?
 - Need to be detectable or contained
- What is implicit coordination?
 - E.g. Listening and observing, randomization
 - No explicit signaling, no dedicated, central controllers

Examples

- TCP way of implicit communication: packet loss
- Local address assignment: Usage of passive duplicate address detection
- WLAN MAC protocols with Carrier Sense Multiple Access
- Distributed resource control in multi-hop networks
 - Listening and overhearing to minimize control messages

WLAN MAC Protocol with Implicit Communication

- How to split resources of a shared (wireless) medium?
- Carrier Sense Multiple Access / Collision Detection, (IEEE 802.11 CSMA / CD)
- „The CSMA/CD protocol functions somewhat like a dinner party in a dark room”
 - Charles Spurgeon
- E.g. A and C send data at same time to B => Collision
 - A and C may not hear each other
- Collision events give information about status of the network
 - Capetanakis, 1979
- Random backoff after collision

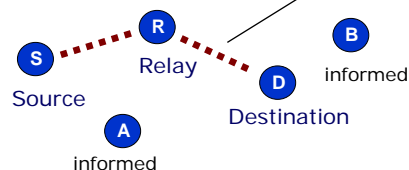


DARE: Distributed Resource Control in Multi-hop Networks

Goal: Extend WLAN 802.11 MAC to make it suitable for

- Resource allocation for real-time applications
- Distributed control of multi-hop reservations

Approach: End-to-end reservation of resources (time slots)

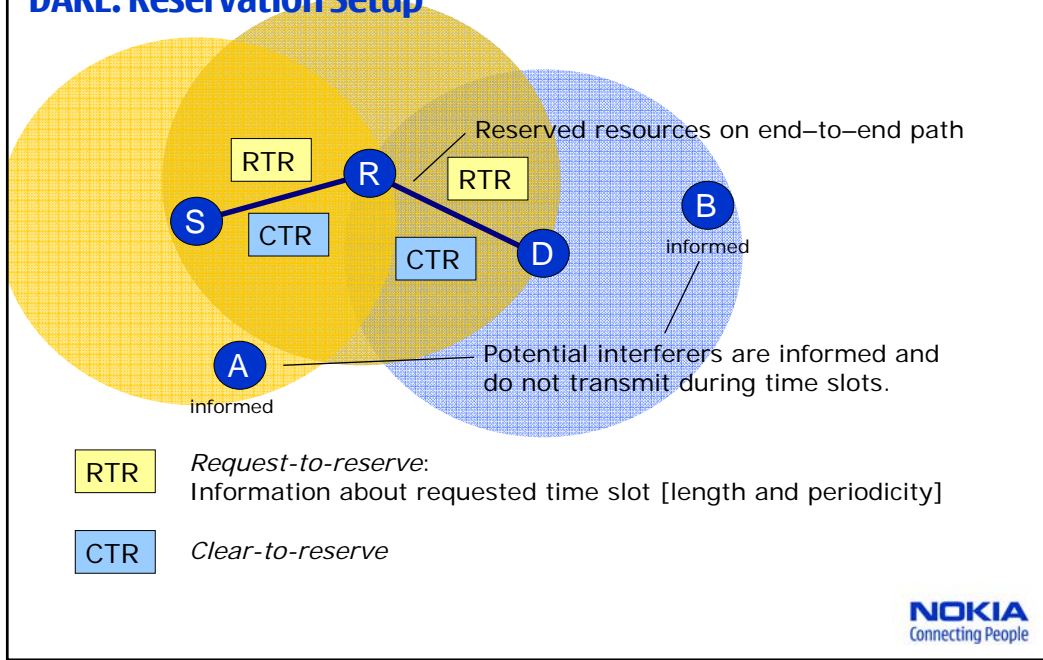


Protocol: DARE

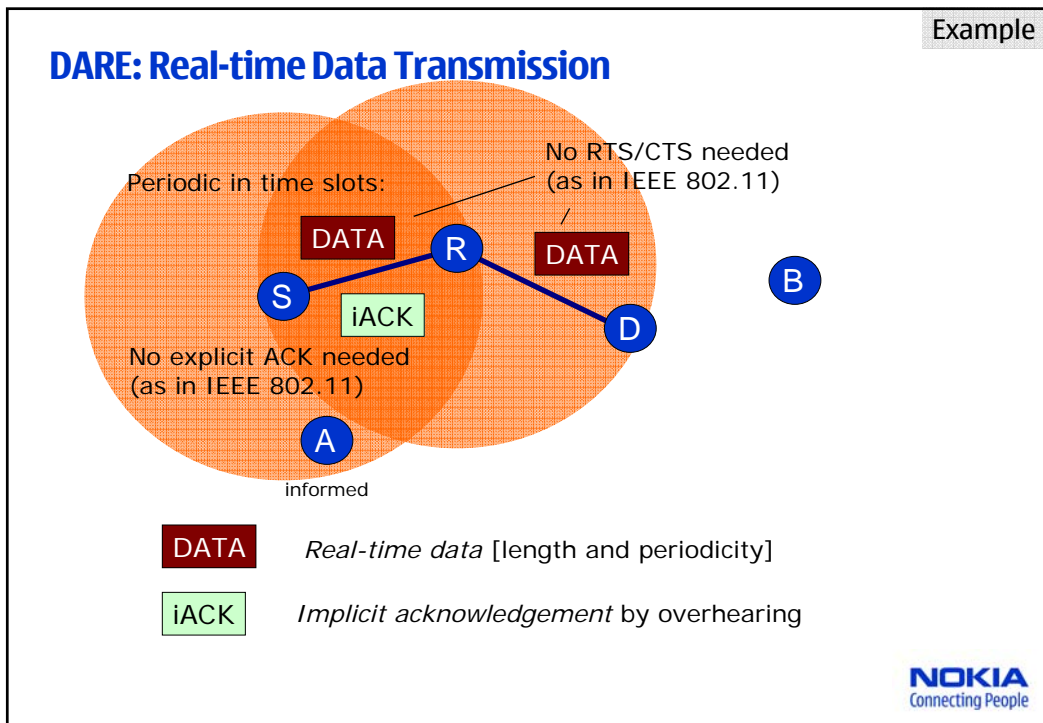
Distributed allocation of time slots for real-time traffic end-to-end

E. Carlson, C. Prehofer, C. Bettstetter, H. Karl, A. Wolisz.
A Distributed End-to-End Reservation Protocol for IEEE 802.11-based Wireless Mesh Networks.
IEEE Journal on Selected Areas in Communications (JSAC), 2006.

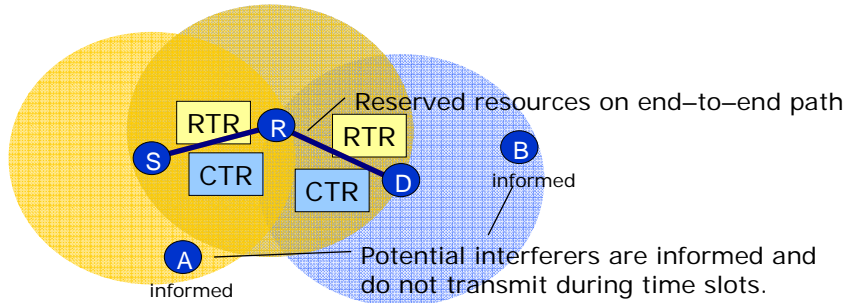
DARE: Reservation Setup



DARE: Real-time Data Transmission



DARE: Protocol Design Summary



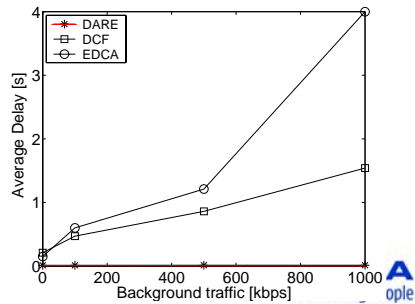
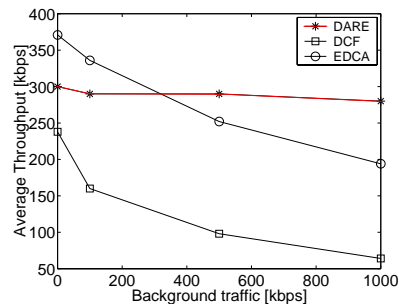
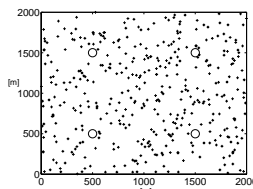
- Locally controlled end-to-end reservations
 - Nodes on reserved path maintain (soft) reservation state
 - Need to inform vicinity about reserved slots
 - Need to agree with other reservations in vicinity – first come first serve basis
- Continuously informing other nodes about reservations
 - „Biggy-packing“ of information on other messages
 - How far to spread information about reserved slots – vs reuse of time slots
 - Especially difficult in case of mobility
- No central coordination or optimization of reservations and time slots

Performance Result: Dare vs EDCA (802.11 with Priorities)

- EDCA to give priority to real-time traffic
 - Retransmission timers, node queues
- DARE gives higher throughput and stable delays for higher load
 - DARE creates minimal signaling overhead

Evaluation Settings with NS2 Simulator

- 400 nodes in 2 x 2 km square, 230m range
- Random on-off of nodes
- IEEE 802.11 DCF and EDCA
- 10 Real-time flows with 40kb/s
- Background traffic with 20 or 50kb/s
- AODV routing protocol

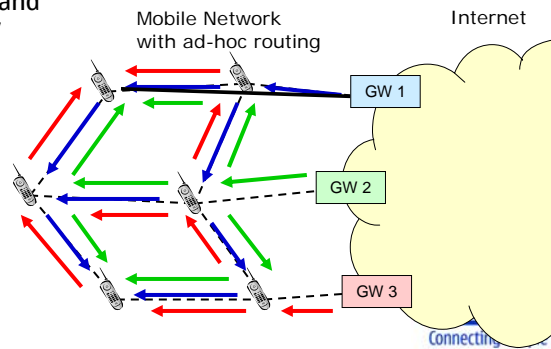


Paradigm #3: Minimise Long-lived State Information

- Reduce state which is **synchronized** between two or more nodes

Example

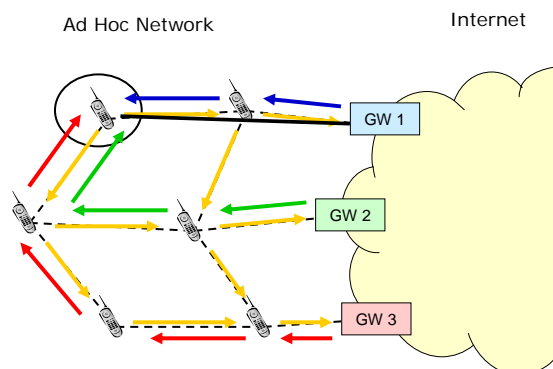
- Discovery of gateways in multi-hop access network
 - Self-organizing ad-hoc network(s) connected to Internet via gateway (GW)
 - Nodes discovery optimal gateway and obtain global internet connectivity
- Comparison of proactive and reactive discovery
- Proactive discovery
 - Gateways send periodic announcement
 - Nodes select best gateway



Reactive Discovery of Gateways

Example

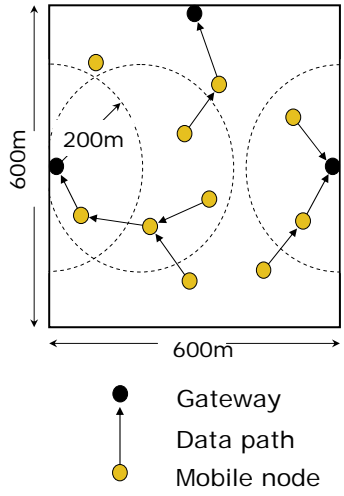
- Reactive Discovery
 - Node sends announcement
 - Gateways reply
 - Nodes select best gateway
- Hybrid discovery option
 - Gateways in proactive mode
 - Nodes only proactive if gateway known



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Gateway Selection & Handover with Mobile Nodes

- MRAN protocol for gateway discovery/selection
 - Autoconfiguration of global address
 - Packet forwarding to/from fixed network
- Gateway selection strategies and handover
 - Proactive and reactive gateway selection
 - **Forced handovers** to gateway if old gateway lost
 - **Optimizing handovers** to closer gateway
- Testbed implementation
 - Protocol implementation in Linux
 - 13 nodes conected via WLAN
 - Mobility emulated (via packet filtering)
 - Different ad-hoc routing protocols (AODV, OLSR)



P. Hoffmann, C. Bettstetter, C. Prehofer,
**Performance Impact of Multihop Handovers in an IP-based
 Multihop Radio Access Network**, ACM SIGMOBILE Mobile Computing
 and Communications (MC2R) Review, 2006.

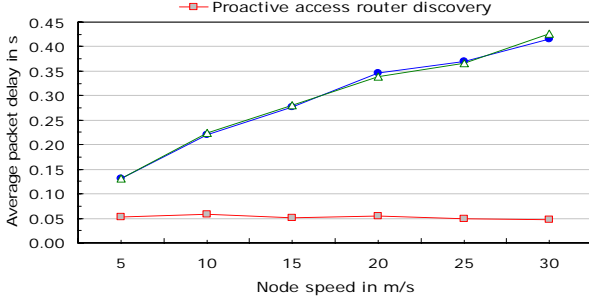
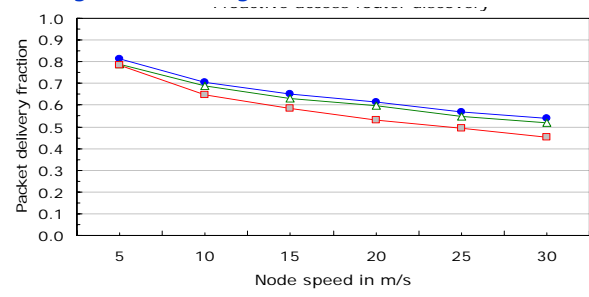
IETF draft submitted 2006
 draft-hofmann-
 autoconf-mran-00.txt **NOKIA**
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Reactive vs Proactive Gateway Discovery

Evaluation Results

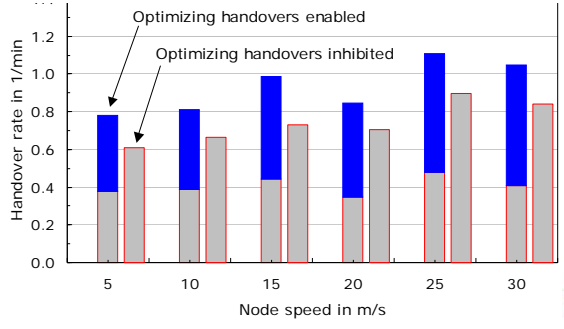
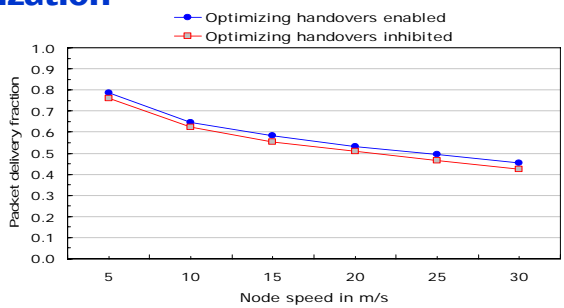
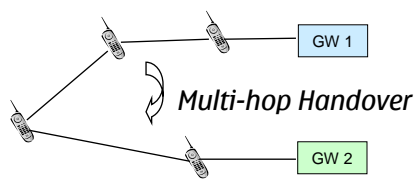
- Little impact on data delivery fraction
- Proactive discovery performs better for higher mobility

- Parameters for performance evaluation
- Number of mobile nodes 10
 - Number of gateways 3
 - Constant bit rate of each MN 10 kbit/s
 - Measurement period 30 min
 - Number of experiments 10

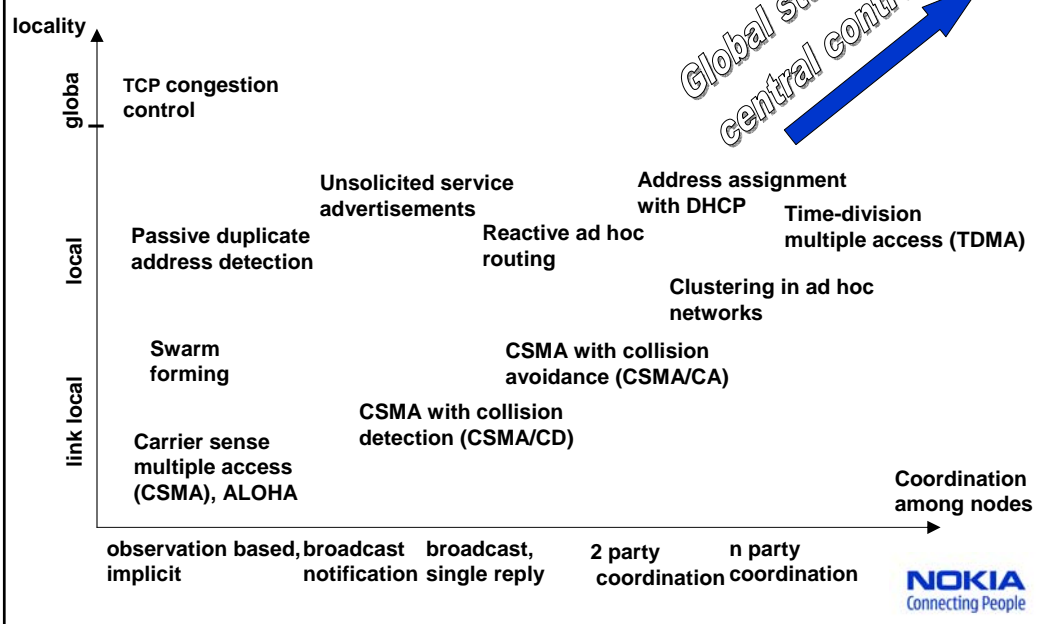


Multi-hop Handover Optimization

- Optimizing handovers faster than forced ones
 - Optimizing ~ 5ms
 - Forced ~ 35 ms
- Optimizing handovers increase number of handovers
- Optimizing handovers improve system performance



Locality and State



Paradigm #4: Adaptation

„Adapt or perish, now as ever, is nature's inexorable imperative“ H. G. Wells

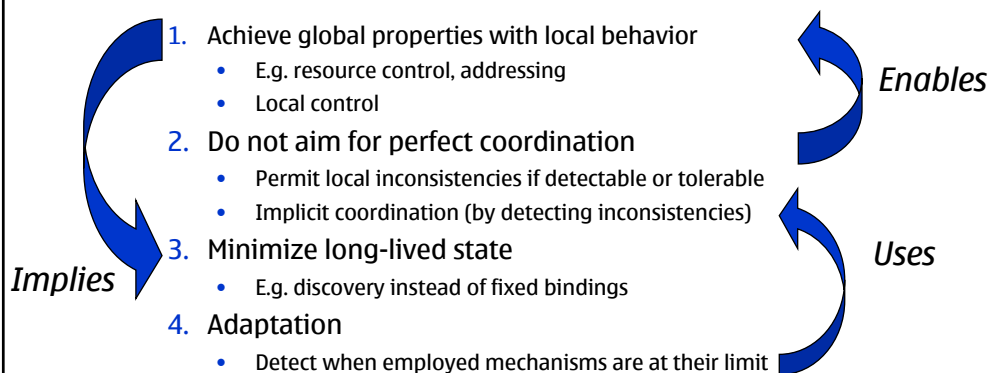
- Need to adapt when local/implicit mechanisms do not suffice
 - Monitor and adapt – often implicit communication

Levels of Adaption

- **Level 1:** A protocol is designed so that it adapts continuously (e.g. failure, mobility)
 - Adapt parameters continuously
 - E.g. TCP transmission window control adapts to round trip time (RTT)
- **Level 2:** A protocol is designed to adapt to specific events and conditions
 - Adapt parameters (e.g. value of timers, cluster size) in order to optimize system performance
 - E.g. TCP reduces sending rate upon congestion
- **Level 3:** A protocol is designed so that it realizes if the changes are so severe that the currently employed mechanism is no longer suitable
 - Major reaction and change of behavior
 - E.g. TCP time out upon consecutive packet loss

Summary of Design Paradigms & Dependencies

What makes self-organization difficult in practice?



From Networking to Pervasive Computing & Services

Example

How to design pervasive applications?

- Where to maintain state information?
- What level of control can we get?
- How to adapt to changing environment?
 - Robustness is key
- Hope to cope with heterogeneous environment?
 - Devices, protocols, radio
 - E.g. short range of RFID tags



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Example: Opportunistic data collection

Example

- Can we use the same principles for pervasive computing?

Local book finder example:

- All books in a building have RFID tags
- Phones in the building record books in vicinity
- Phones exchange data when people move
 - Optimized data storage and data exchange with network coding
 - Requests can be sent to local devices, e.g. "Find a Java book close by"



Elena Fasolo, Christian Prehofer, Michele Rossi, Qing Wei, Jörg Widmer, Andrea Zanella and Michele Zorzi,
Challenges and New Approaches for Efficient Data Gathering and Dissemination in Pervasive Wireless Networks, Intersense 2006

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Summary and Outlook

Design principles of self-organization

- Identification of main paradigms of self-organized networks
- Self-organization leads to new research approaches and system design
- Towards „design patterns“ and methodology for self-organized systems

Related trends & opportunities

- Context-aware pervasive systems
 - How to manage context information?
- Web services & service oriented architecture for pervasive systems
 - Stateless services and „document-oriented computing“
- Self-organized security mechanisms
 - Bio-inspired anti-body generation for mal-ware detection